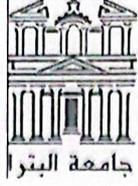


University of Petra		 جامعة البترا - ثلاثون عاما University of Petra
Faculty of Information Technology		كلية تكنولوجيا المعلومات
Department of Computer Science		قسم علم الحاسوب

Advanced Algorithms
601326
Midterm Exam – 2024 2

Instructions for the Exam:

- Write your name and ID number on the exam and answer sheets.
- Write the number of the section that you enrolled in.
- Write the name of your instructor.
- Questions in the exam not allowed.
- Using any type of technology (mobiles, smart watches) not allowed
- Using extra papers or sheets not allowed
- The exam consists of Six questions.

For instructor use only:

Question number	Course ILO	Program ILO	Question weight	Student mark
Q1			4	0.5
Q2	K2		4	4
Q3	I1		7	6
Q4			5	5
Q5			6	2
Q6			4	3 10 (4)
Total /30				20.5

Q1) Prove by induction that:

(4 marks)

$$\sum_{i=0}^n 2^i = 2^{n+1} - 1$$

(0,5)

- let $n=1$

$$\sum_{i=0}^1 2^i = 2^2 - 1 \Rightarrow 1 = 3$$

?

Q2) Consider the following algorithm.

(4 marks)

Algorithm XYZ ($A[0..n-1]$)

//Input: An array $A[0..n-1]$ of n real numbers

$val \leftarrow 5$

$res1 \leftarrow 0$

$res2 \leftarrow 1$

for $i \leftarrow 0$ to $n-1$ do

if $A[i] \leq val$

$res1 \leftarrow A[i] + res1$

if $A[i] > val$

$res2 \leftarrow A[i] * res2$

return $res1 - res2$



- What does this algorithm compute?
- What is the time complexity for the algorithm?

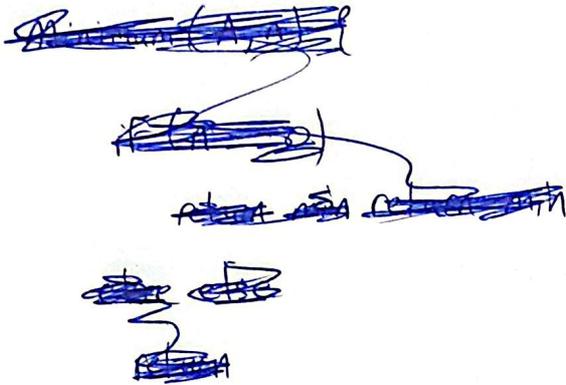
a) the algorithm return $\left(\begin{array}{c} \text{جمع الأعداد} \\ 5 \text{ الأقل من } 5 \end{array} \right) - \left(\begin{array}{c} \text{حزب الأعداد} \\ 5 \text{ الأكبر من } 5 \end{array} \right)$

b) $5n+1$ $O(n)$

Not decrease & conquer

Q3) Design a divide-and-conquer algorithm to find Minimum Number in an unsorted array of N elements (without using any of the sorting algorithms covered in the class). Setup a recurrence relation for your algorithm. (7 marks)

6



A array

~~A array~~

~~A array~~

```
minimum(A, n, min) {  
    if (n == 0)  
        return min;  
    else if (A[n] < min)  
        min = A[n];  
    return minimum(A, n-1, min);  
}
```

$$T(n) = \begin{cases} 0 & , n = 0 \\ T(n-1) & , n > 0 \end{cases}$$

Q4) Consider the following recursive algorithm:

(5 marks)

5

```
Algorithm Q(n)
//Input: A positive integer n
if n = 1
    return 1
else
    return Q(n - 1) + 2 * n - 1
```

- Setup a recurrence relation for the number of multiplications made by this algorithm.
- Solve the recurrence relation using backward substitution and find Big O.

$$a) T(n) = \begin{cases} 0 & , n = 1 \\ T(n-1) + 1 & , n > 1 \end{cases}$$

$$b) T(n) = T(n-1) + 1$$

$$T(n-1) = T(n-2) + 1$$

$$T(n-2) = T(n-3) + 1$$

$$T(n-3) = T(n-4) + 1$$

$$T(n) = T(n-4) + 4$$

$$T(n) = T(n-i) + i$$

$$n-i = 1 \Rightarrow i = n-1$$

~~n-i = 1~~

$$T(n) = T(n-n+1) + n-1$$

$$T(n) = n-1$$

$$O(n)$$

Q5) Given the following array, $A = -3, 8, 4, 22, 13, 6, 20, 9$

(6 marks)

②

- Apply Quick Sort algorithm to sort the array elements in ascending order (show detailed steps)
- Compare Quick sort to Insertion Sort in terms of best and worst cases.

a) ~~using method of j for the pivot~~

~~partition~~ $S \leftarrow \text{partition}(A[L \dots r])$

pivot

~~partition~~ = -3

~~i = 1~~

~~j = 9~~

repeat — repeat $i = i + 1$ until $A[i] > \text{pivot}$
 repeat $j = j - 1$ until $A[j] < \text{pivot}$
 swap $A[i], A[j]$

until $i \geq j$

swap $A[i], A[i]$

swap pivot, $A[j]$

return j

-3 8 4 22 13 6 20 9

2 4 22 13 6 20 9

~~4 22 13 6~~

But you should the Quick Sort algo. For sorting the array NOT to write the algo. itself !!

b) - using the above quick sort will be in the worst case

②

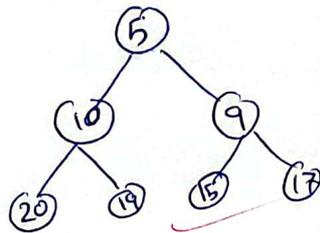
because ~~every time we select a pivot~~ the pivot always sorting at the left $(n-1)$ every time — time complexity will be ~~worst case~~ if the pivot was in the middle we will get $O(n^2)$ a best case $O(n \log n)$ — only if it divide the algorithm into two equal parts.

- insertion sort best case when its already sorted $O(n)$

worst case when its reversed sorted $O(n^2)$

Q6) Show that the worst case running time of MAX-HEAPIFY on a heap of size n is $\log n$
(Hint: For a heap with n nodes, give node values that cause MAX-HEAPIFY to be called recursively at every node on a simple path from the root down to a leaf.) (4 marks)

3.42



the worst case will occur when its reversed sorted
building max heap will call heapify everytime

~~everytime~~

?